
UniLFS: A Unifying Logical Framework for Service Modeling and Contracting

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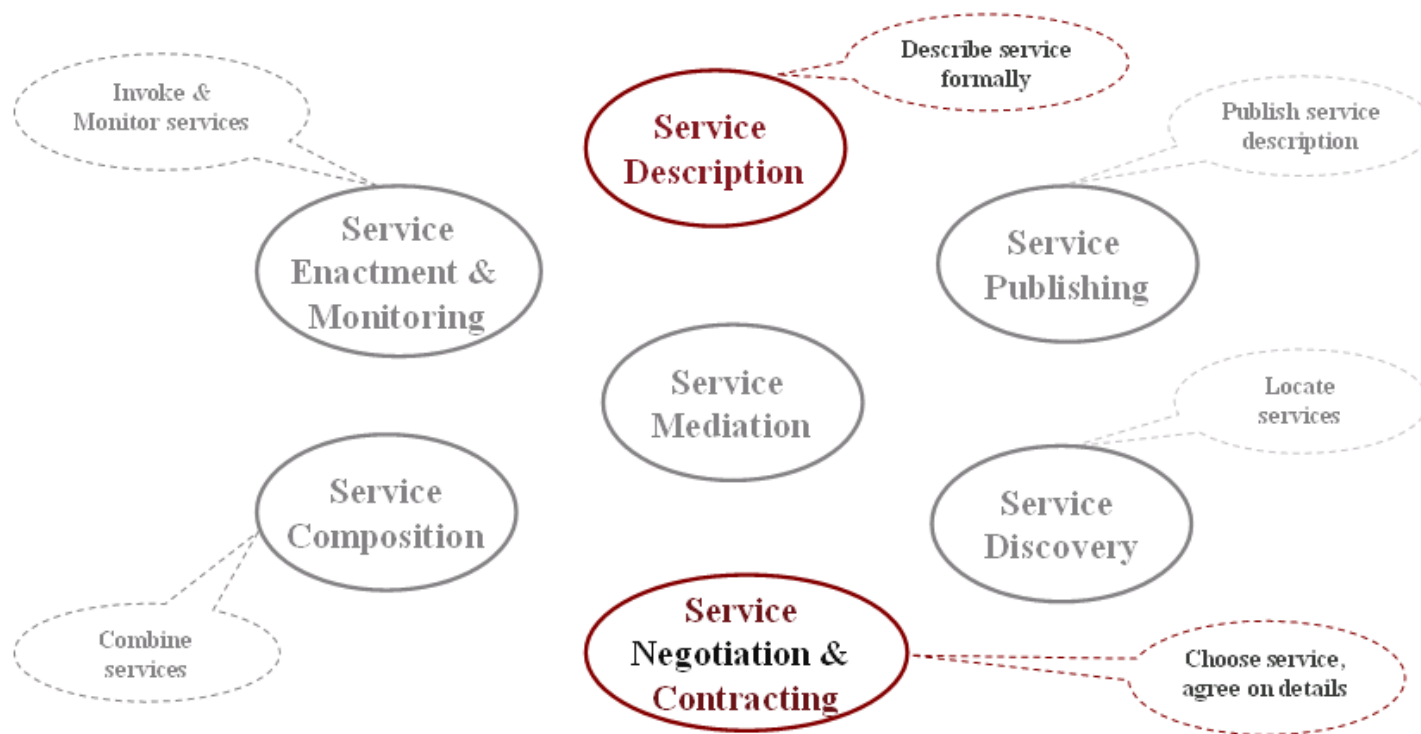
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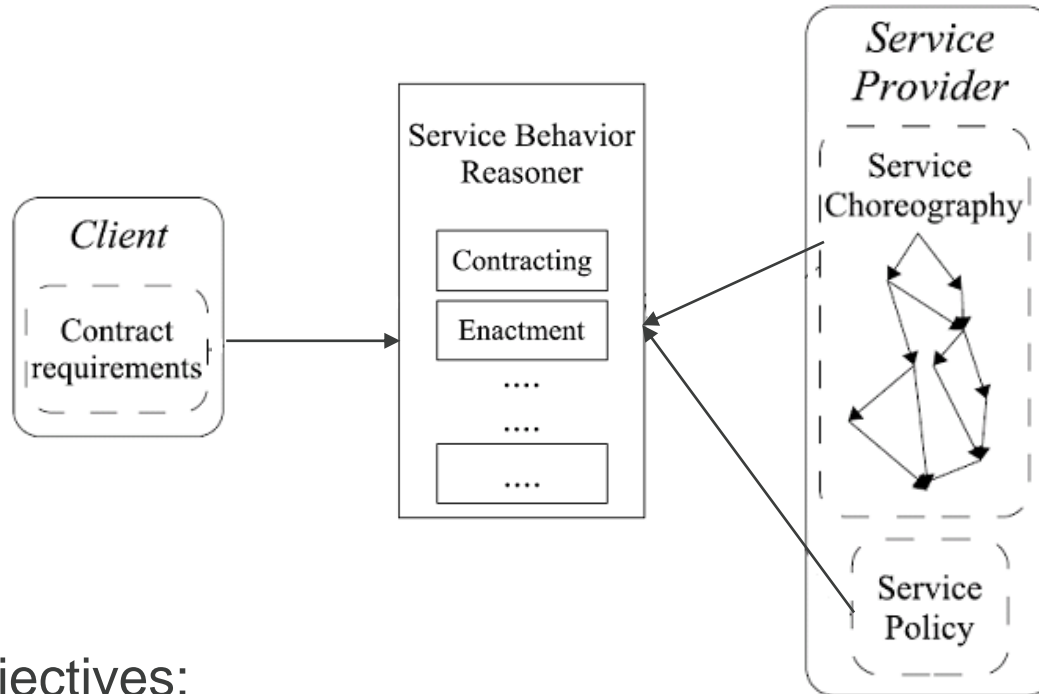
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- Context and motivation: Automated service contracting
- **UniLFS: Modeling** and **reasoning** about service behavior for automated service contracting
 - Elements of service contracts
 - Expressivity of the modeling language for service contracts
 - Reasoning technique for service contracting
- Related work
- Conclusions and outlook



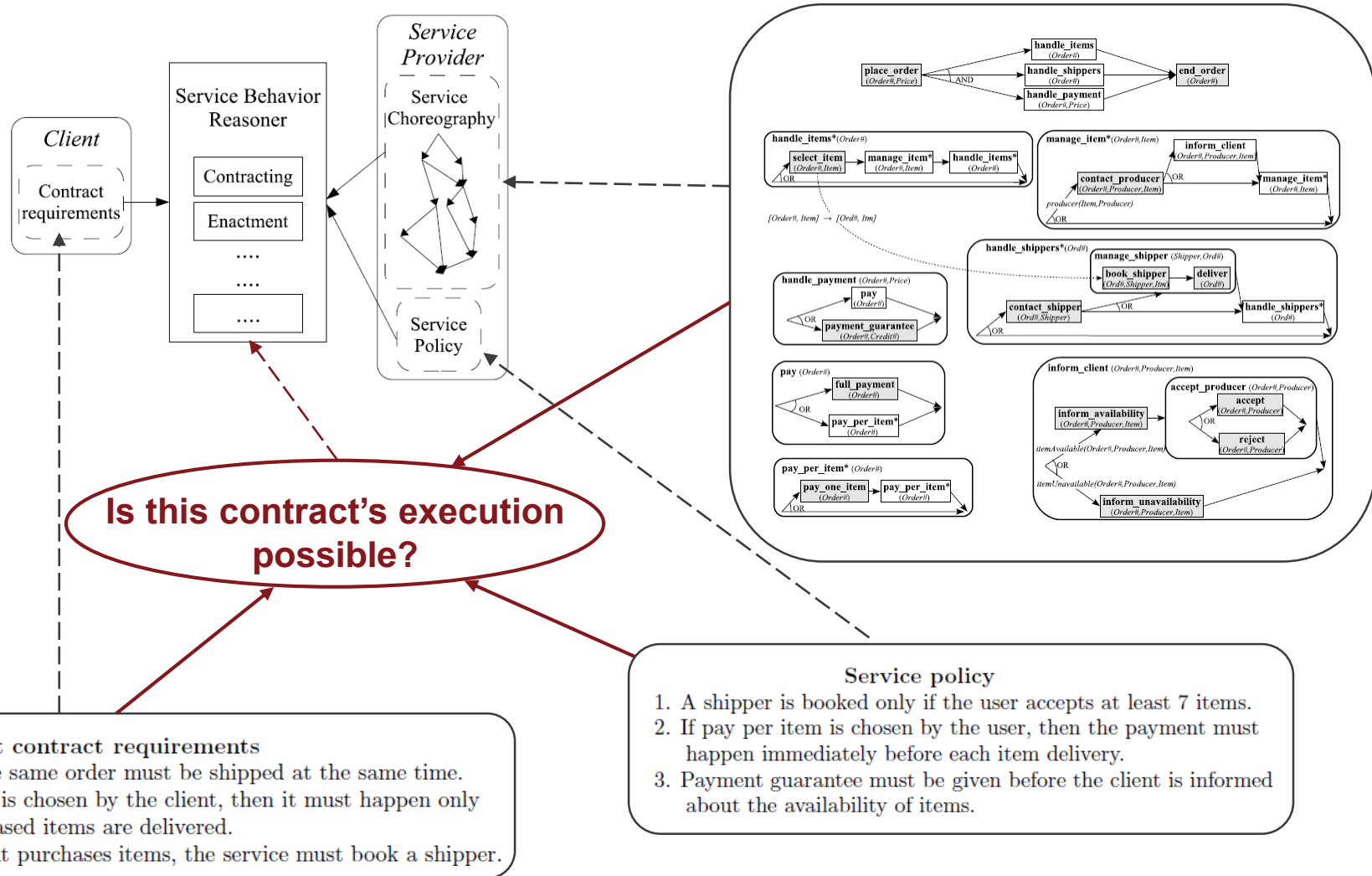
- **Overall aim:** Automate service contracting
 - Problems with traditional service contracting
 - High costs for contract establishment and management
 - Slow for rapidly changing business situations
 - Automated service contracting can potentially
 - Lower contract establishment and management costs
 - Speed up and improve contract establishment and management
- **Current status:** Limited expressivity of modeling frameworks and lack of sound and complete contract establishment techniques
- **General approach**
 1. Focus on those aspects of contracting that constrain interactions in processes
 2. Define the conceptual framework
 3. Formalization, representation, and reasoning

Service Contract = Choreography + Service Policies + Client Contract Requirements



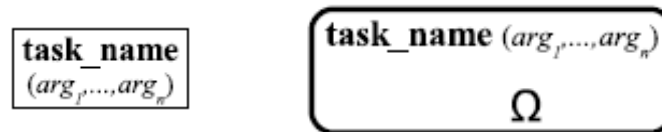
Main objectives:

1. An expressive **language** for contracting
 - Service choreography
 - Service policies
 - Client contract requirements
2. A **reasoning mechanism** to decide if service contracting is possible

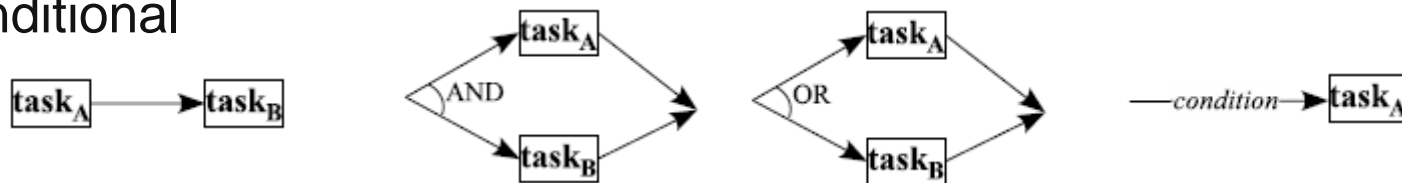


- Aim
 - An expressive representational framework for service contracts: combining procedural and declarative aspects
- Common requirements
 - Service choreographies
 - Conditional control flow: Sequential, concurrent, non-deterministic, and iterative interactions
 - Data flow: Local and non-local data passing between interactions
 - Service Policies and Client Contract Requirements
 - Temporal, conditional, and data constraints

- Tasks: Task atoms and task definitions (hierarchical representations)



- Composition of tasks: Sequential, concurrent, non-deterministic, and conditional

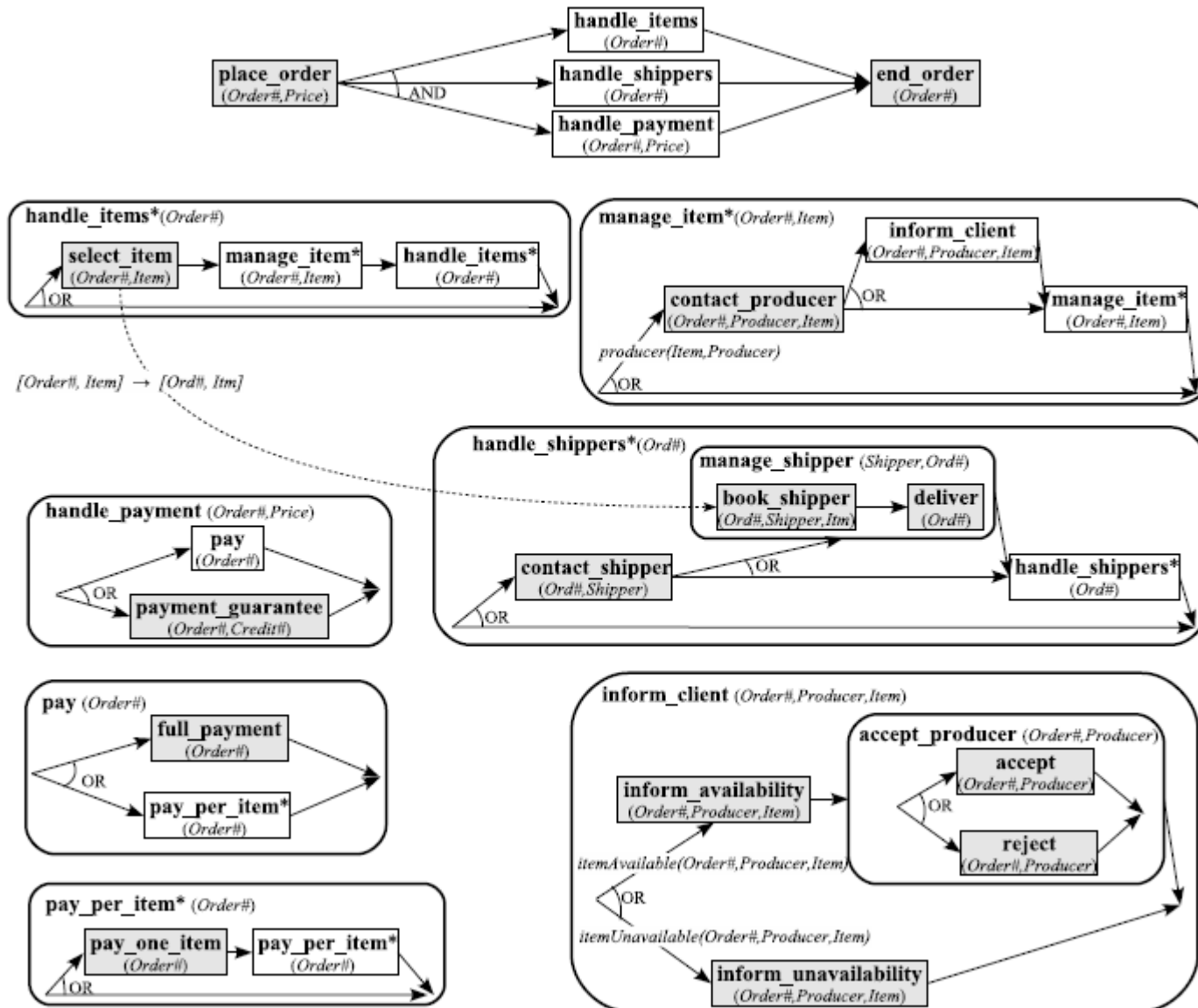


- Data passing between tasks

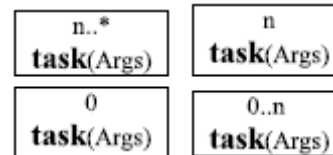
- Data-passing through **shared arguments** is possible between a task and its direct successors, or within the definition of a composite task
- Data-passing through **shared data space** is used when passing data is not possible through shared arguments due to the difference in scope of the arguments



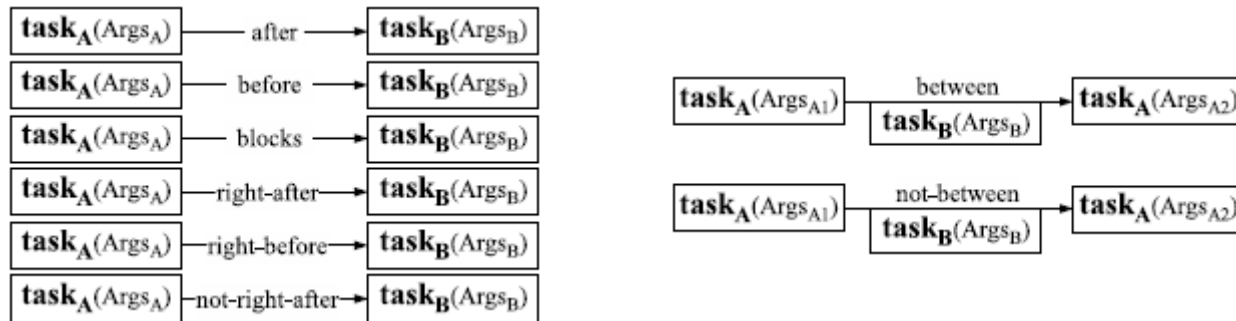
Example



- Existence constraints

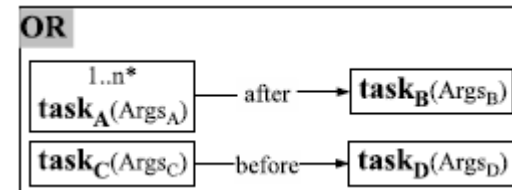
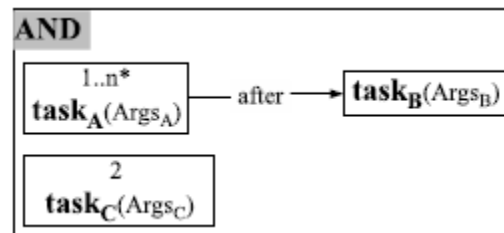


- Serial constraints



- Complex constraints – if C1 and C2 are constraints, then so are

- C1 **and** C2
- C1 **or** C2



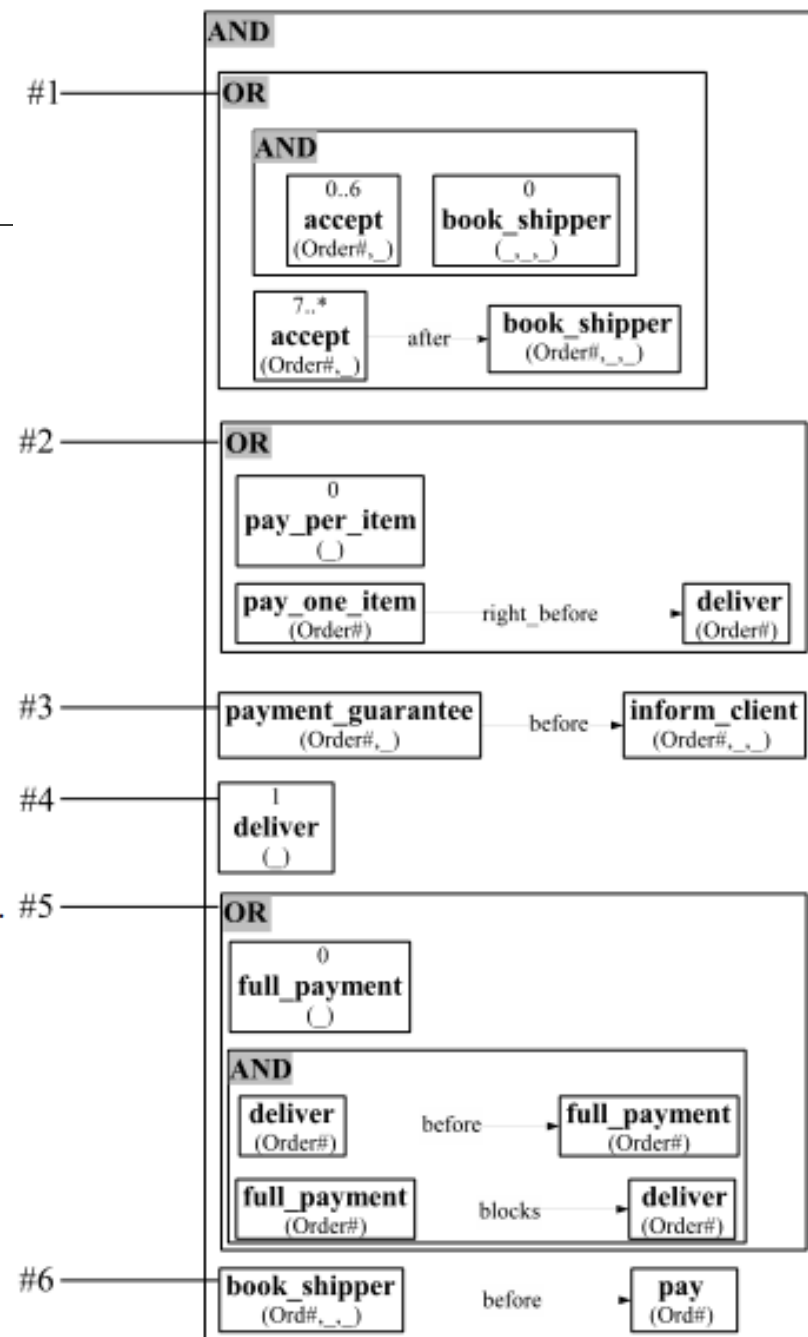
Example

Service policy

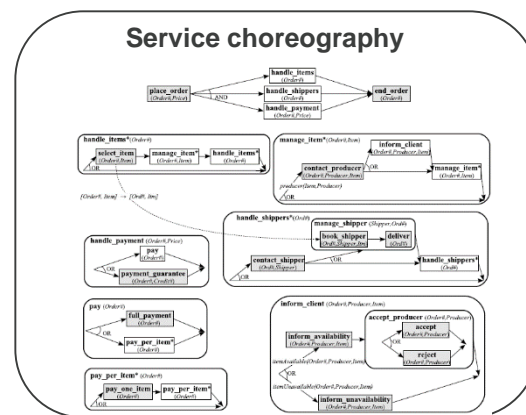
1. A shipper is booked only if the user accepts at least 7 items.
2. If pay per item is chosen by the user, then the payment must happen immediately before each item delivery.
3. Payment guarantee must be given before the client is informed about the availability of items.

Client contract requirements

4. All items in the same order must be shipped at the same time.
5. If full payment is chosen by the client, then it must happen only after all purchased items are delivered.
6. Before the client purchases items, the service must book a shipper.



- **Service contracting:** **Decide** if the contract is satisfiable
- **Service enactment:** If it is, **find an enactment**



What is the actual order of interactions if service contracting is possible?

Client contract requirements

4. All items in the same order must be shipped at the same time.
5. If full payment is chosen by the client, then it must happen only after all purchased items are delivered.
6. Before the client purchases items, the service must book a shipper.

Service policy

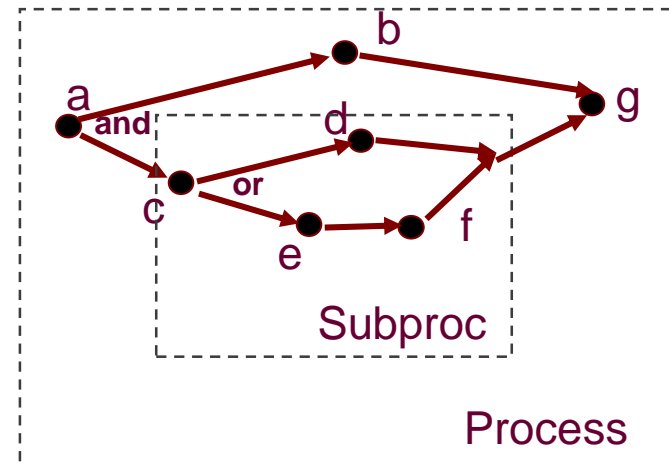
1. A shipper is booked only if the user accepts at least 7 items.
2. If pay per item is chosen by the user, then the payment must happen immediately before each item delivery.
3. Payment guarantee must be given before the client is informed about the availability of items.

1. Concurrent Transaction Logic (CTR) – our formalism for
 - Choreographies
 - Service policies
 - Client contract requirements
2. Extend CTR when necessary
3. Prove correctness

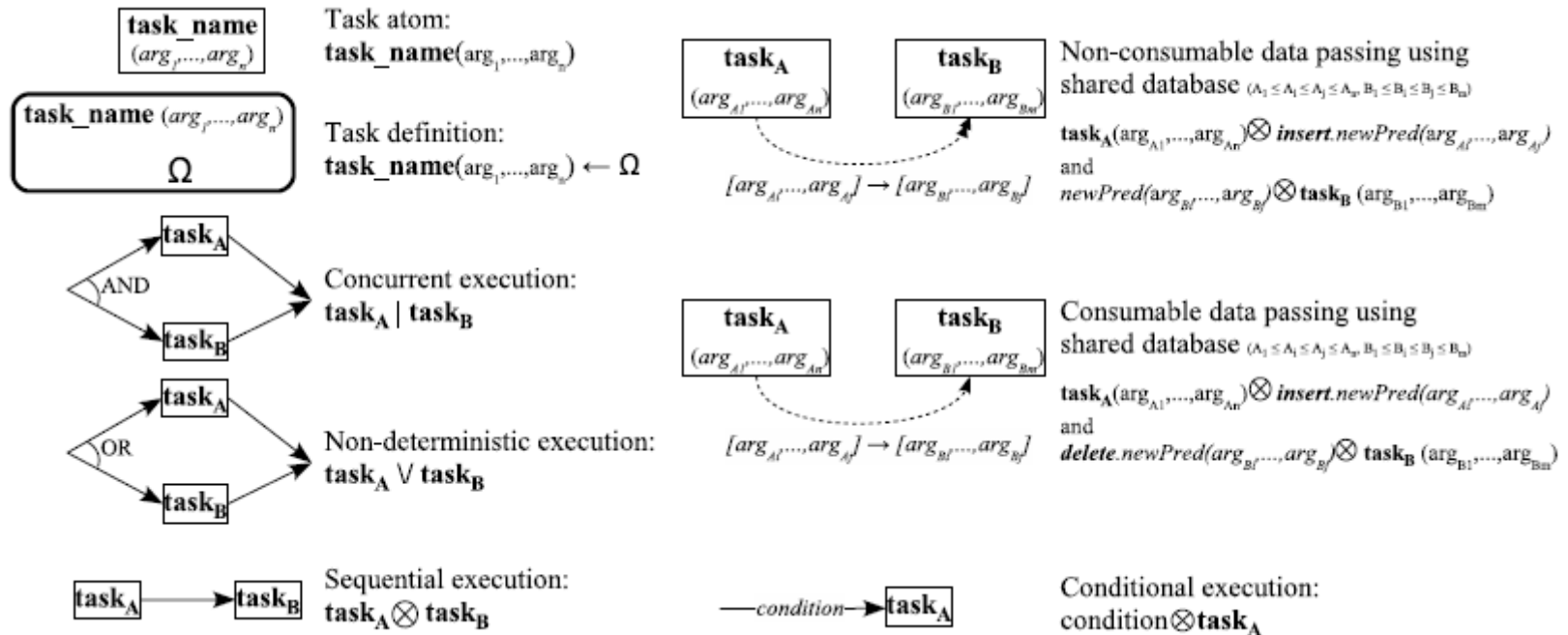
- Why CTR?
 - Logic for specification and execution of transactional processes
 - Integrates queries, updates, and transaction composition
 - Has proof procedure for concurrent Horn formulas: *proof = execution*
 - Can express a wide variety of constraints
 - Has been applied in the area of workflow verification and scheduling with promising results
- Atomic CTR formulas: same as in classical logic
- More complex formulas: built using connectives $a \otimes b$, $a \mid b$, $a \wedge b$, $a \vee b$, $\neg a$, $\odot a$

$Process \leftarrow a \otimes (b \mid Subproc) \otimes g$

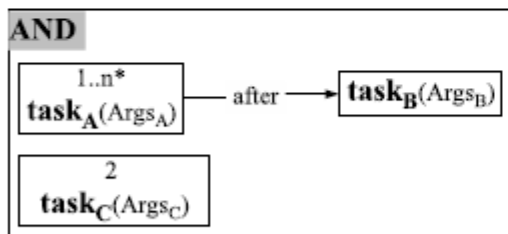
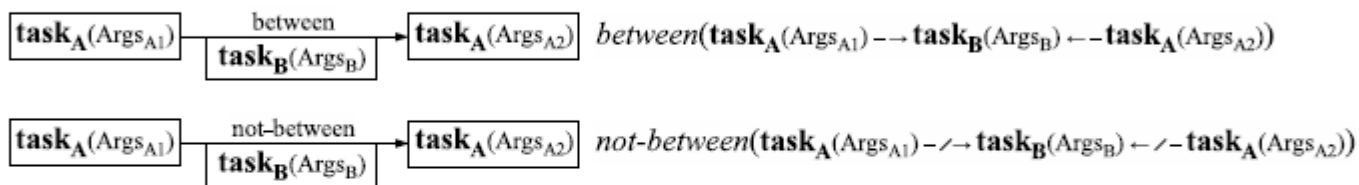
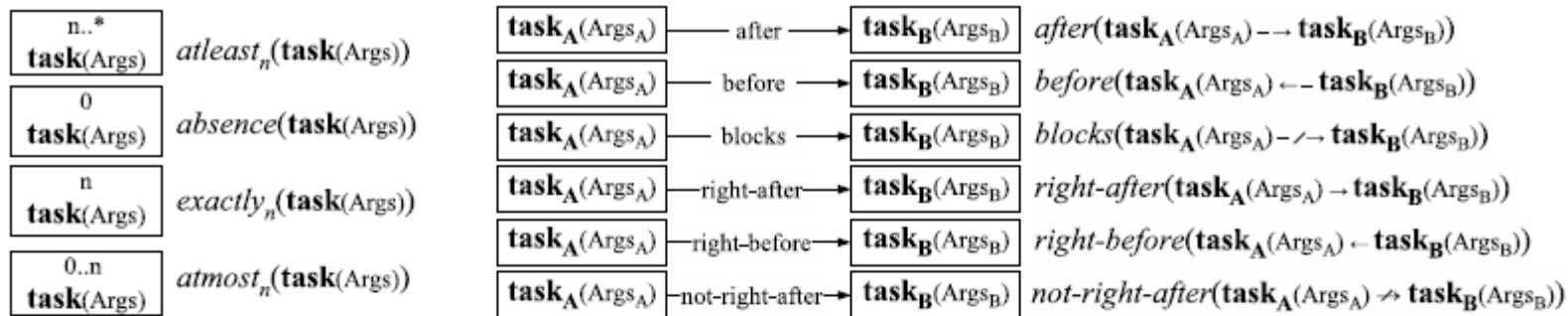
$Subproc \leftarrow (c \otimes (d \vee (e \otimes f)))$



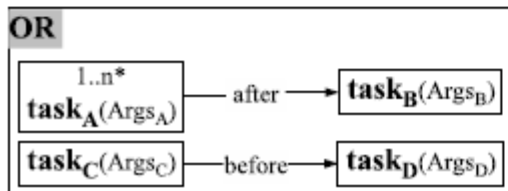
Graphical representation of a service choreography in UniLFS



Graphical notation for constraints in UniLFS



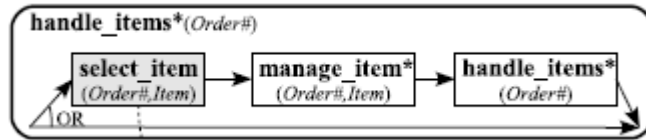
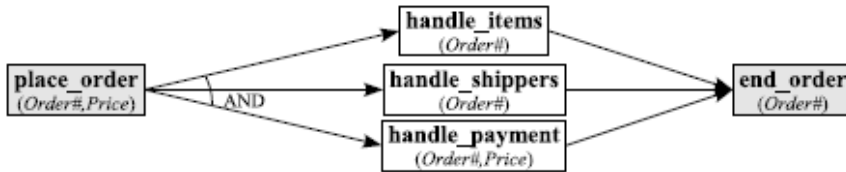
Conjunctive composition
 $atleast_1(\mathbf{task}_A(\mathit{Args}_A)) \wedge after(\mathbf{task}_A(\mathit{Args}_A) \dashrightarrow \mathbf{task}_B(\mathit{Args}_B)) \wedge exactly_2(\mathbf{task}_C(\mathit{Args}_C))$



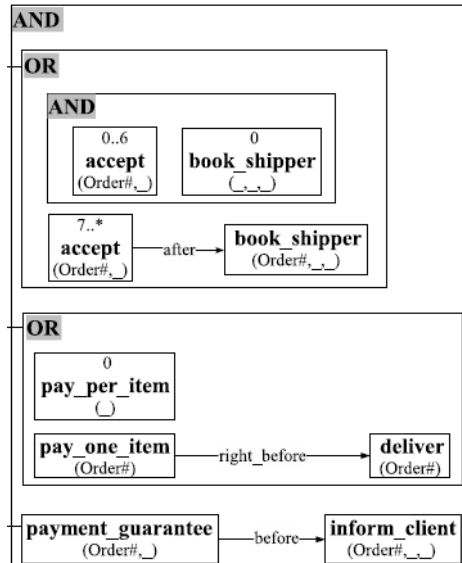
Disjunctive composition
 $(atleast_1(\mathbf{task}_A(\mathit{Args}_A)) \wedge after(\mathbf{task}_A(\mathit{Args}_A) \dashrightarrow \mathbf{task}_B(\mathit{Args}_B))) \vee before(\mathbf{task}_D(\mathit{Args}_D) \leftarrow \mathbf{task}_C(\mathit{Args}_C))$

More examples of constraints that can be captured in UniLFS

- if p is executed, then q must also execute (before or after q)
- if p is executed, then p must also be executed, and vice versa
- every occurrence of task p must be followed by an occurrence of task q with the same argument and there must be an occurrence of p before every occurrence of q and their arguments must be the same
- if task p is executed then q must execute after it, and before that q there can be no other p
- if task q is executed, it has to be preceded by an occurrence of p ; the next instance of q can execute only after another occurrence p
- tasks p and q must alternate
- ...



...



...

ConcurrentHorn \wedge Constraints

```
place_order(?Order#, ?Price)  $\otimes$ 
(handle_items(?Order#)
 | handle_shippers(?Order#)
 | handle_payment(?Order#, ?Price))  $\otimes$ 
end_order(?Order#)
```

```
handle_items(?Order#)  $\leftarrow$ 
(select_item(?Order#, ?Item)  $\otimes$ 
 insert.selected_item(?Order#, ?Item)  $\otimes$ 
 manage_item(?Order#, ?Item)  $\otimes$ 
 handle_items(?Order#))
handle_items(?Order#)  $\leftarrow$  state
```

...

\wedge

1. (atmost₆(accept(_Order#, _)) \wedge absence(book_shipper(_, _))) \vee (atleast₇(accept(_Order#, _)) \wedge after(accept(_Order#, _) \rightarrow book_shipper(_Order#, _))) where atmost_n(p) is a shorthand for absence(p) \vee exactly₁(p) \vee exactly₂(p) \vee ... \vee exactly_n(p)
2. absence(pay_per_item(_)) \vee right_before(pay_one_item(_Order#) \leftarrow deliver(_Order#))
3. before(inform_client(_Order#, _) \leftarrow payment_guarantee(_Order#, _))

...

- **Prior work**

- Original CTR proof theory does not handle constraints
- Prior work added simple constraints, but
 - No iterative processes
 - No constraints on iterations
 - No data flow

- **New**

- Add the necessary modeling primitives
- Extend the CTR proof theory to handle service contracting with iterations, data flow, constraints, etc.

- **As a result**

- Can capture much of BPMN and WS-BPEL process modeling, extend them with declarative constraints
- And contracting over them

- **Contracting**: Find out if an execution of the CTR formula $Choreography \wedge Policy \wedge ClientRequirements$ exists
- **Enactment**: Find a constructive proof for this execution

Main result: The extended inference system is **sound** and **complete** for proving the above

- **Soundness**: Every found execution is correct
- **Completeness**: All possible executions can be found

- Sequent: $P, s \dashv\vdash (\exists) \psi \wedge C$
 - Meaning: $(\exists) \psi$ can execute starting at state s given P in such a way that C is satisfied

- Each inference rule looks like this: $\frac{\textit{sequent}_1}{\textit{sequent}_2}$

For example:

$$\frac{P, s \dashv\vdash (\exists) \psi' \sigma \wedge C}{P, s \dashv\vdash (\exists) \psi \wedge C}$$

- Has seven inference rules and one axiom

- E-contracting
 - Pieces of frameworks, models, architectures and different aspects and phases of e-contracting (negotiation, enforcement, violation detection, monitoring, legal aspects), but no unifying formalism
 - We provide a unifying and general framework for e-contracting
- Workflow/process modeling
 - Most languages are procedural (WS-BPEL, BPMN, YAWL).
The declarative ones (DecSerFlow) are inexpressive
 - Our framework: very expressive. Integrates conditional control flow, data flow, hierarchical modeling, and complex constraints
- Process verification
 - Most of the existing approaches use temporal logic/model checking. No obvious way to handle certain aspects: data flow, hierarchical modeling
 - We use CTR. More natural modeling language. Sometimes has better complexity.

- UniLFS summary
 - Identified main element of service contracts and formally formulated the problem of service contracting
 - *Modeling*: Developed an expressive representational framework for service contracts, combining procedural and declarative elements
 - *Reasoning*: Developed a sound and complete proof theory for deciding if contracting for a service is possible
- Future work
 - Semi-automated extraction of UniLFS specifications from text (e.g. contracts)
 - Complexity study, e.g. subsets of constraints for which the verification problem has a better complexity
 - Implementation

Thank you!
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